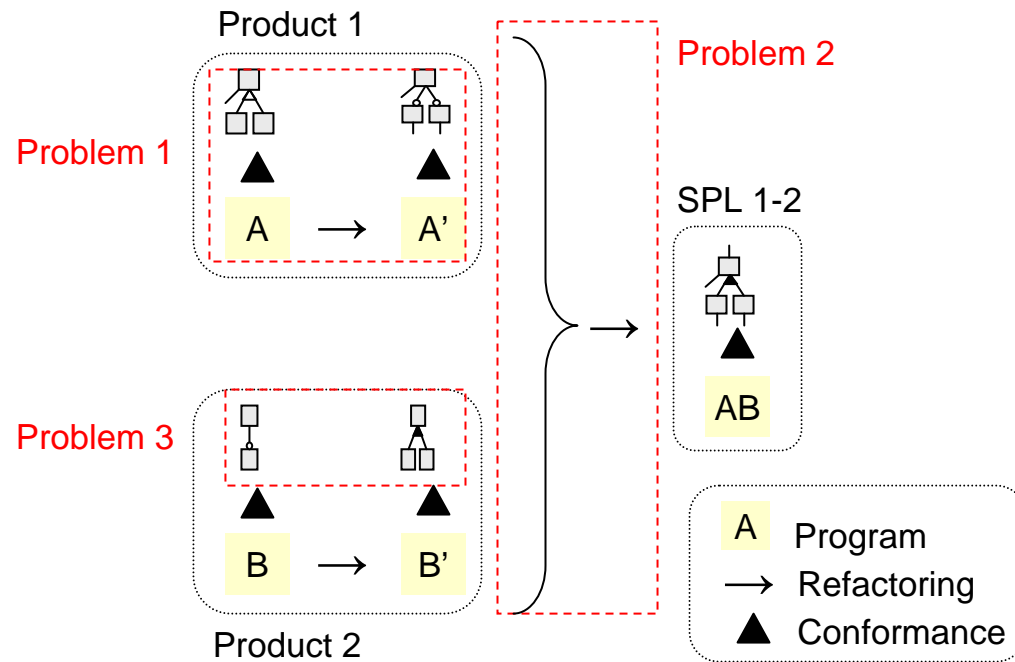

Refactoring Product Lines

Problem

Problems in Refactoring SPL



Feature Model Theory

Abstract Syntax

Name: TYPE

FM: TYPE = [#

features: set[Name],

formulae: set[Formula]

#]

Config: TYPE = [#

config: set[Name]

#]

Formulae

- True
- False
- Name
- Not
- And
- Implies

Dynamic Semantics

```
semantics(fm: {f:FM | wf(f)}): set[Config] =  
  { c: Config |  
    satImpConst(fm,c)  $\wedge$   
    satExpConst(fm,c)  
  }
```

Constraints

satImpConst(fm:{f:FM | wf(f)}, c:Config): bool =
config(c) \subseteq features(fm)

satExpConst(fm:{f:FM | wf(f)}, c:Config): bool =
 \forall f: formulae(fm) | satFormula(f,c)

satFormula(f,c)

True, False : **true, false**

Name(n) : $n \in \text{config}(c)$

Not(f1) : **not** satFormula(f1,c)

And(f1,f2) : $\text{satFormula}(f1,c) \wedge$
 $\text{satFormula}(f2,c)$

Implies(f1,f2) : $\text{satFormula}(f1,c) \Rightarrow$
 $\text{satFormula}(f2,c)$

Static Semantics

wf(fm FM): bool =

$\forall f: \text{formulae}(\text{fm}) \mid \text{wellTyped}(\text{fm}, f)$

wellTyped(fm,f)

True, False : **true**

Name(n) : $n \in \text{features}(fm)$

Not(f1) : wellTyped(fm,f1)

And(f1,f2) : wellTyped(fm,f1) \wedge
wellTyped(fm,f2)

Implies(f1,f2) : wellTyped(fm,f1) \wedge
wellTyped(fm,f2)

Reactive PL Refinement

refines(abs, con:FM): bool =
semantics(abs) \subseteq semantics(con)

Extractive PL Refinement

```
refines(abs1, abs2, con:FM): bool =  
  refines(abs1,con)  $\wedge$   
  refines(abs2,con)
```

Equivalences between Feature Models

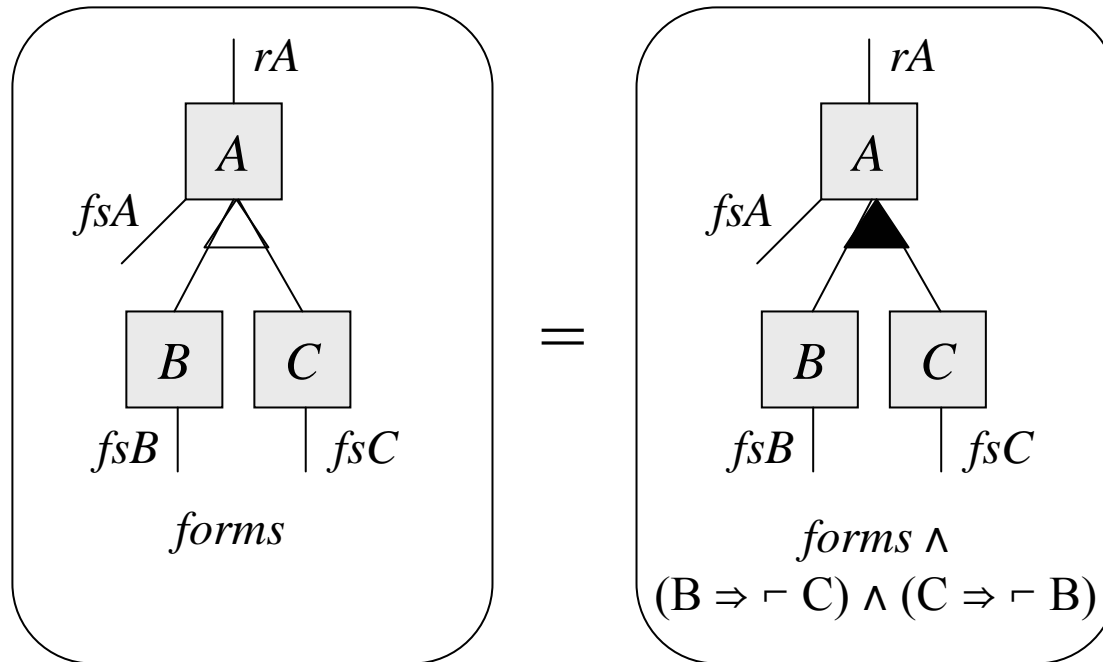
Core Language

- Our core language contains
 - features
 - formulae
- This core language is specified in PVS
- Notice that our core language does not have the same static semantics of the original one.
 - So, some of the syntactic sugar (Laws 4 and 5) will not be used in practice.
 - However, it is as **expressive** as the original one.

Syntactic Sugar Transformations

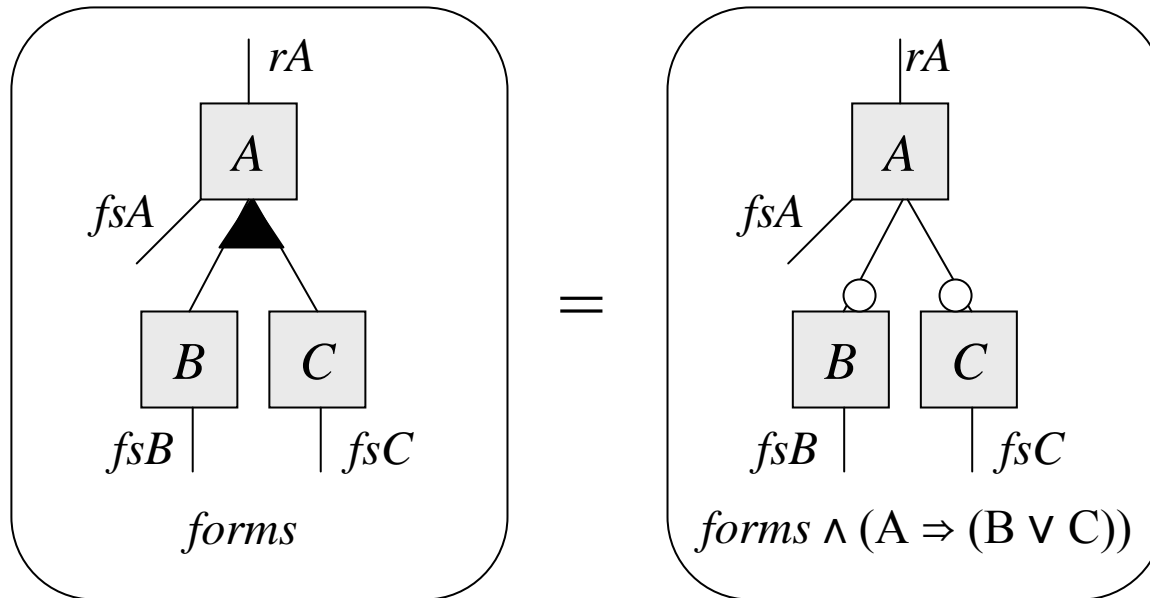
- They are intuitively valid by definition
- We do not need to prove them
- We use the syntactic sugar laws to reduce any feature model to our core language
- We propose 14 syntactic sugar refinements
- This set is sound, minimal and complete.

Replace Alternative (Law 1)

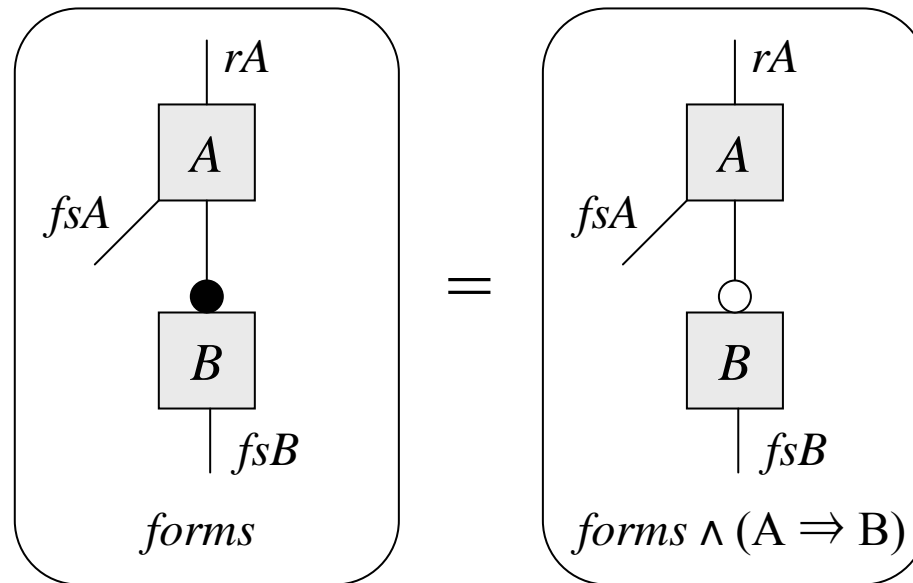


r : root
 fs : features
 $forms$: formulae

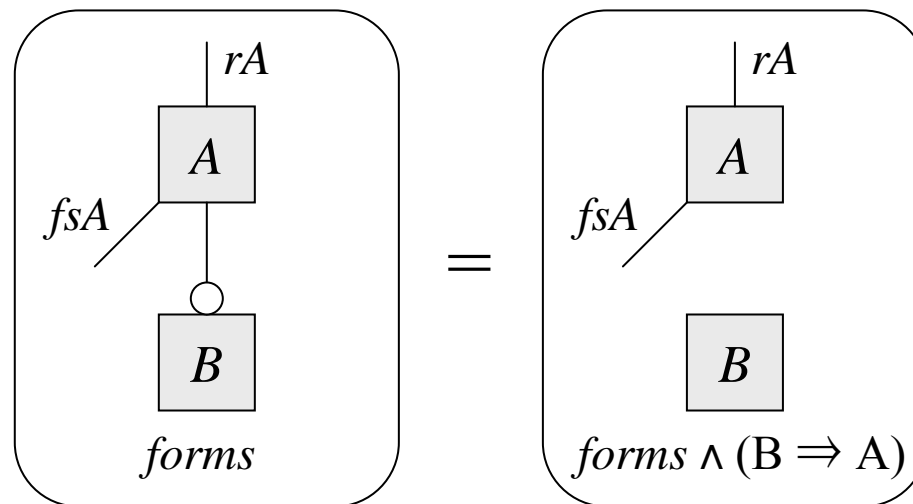
Replace Or (Law 2)



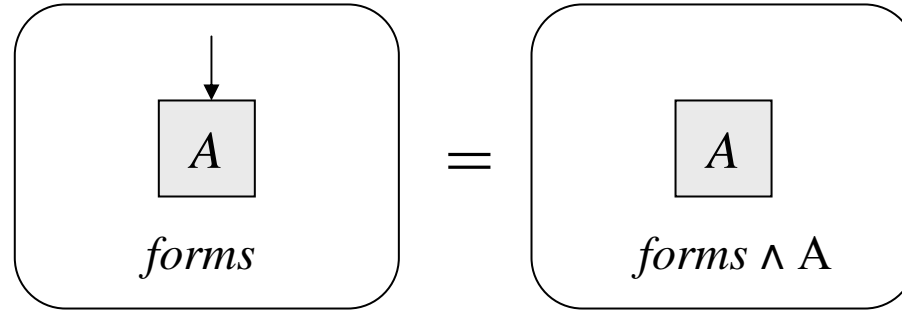
Replace Mandatory (Law 3)



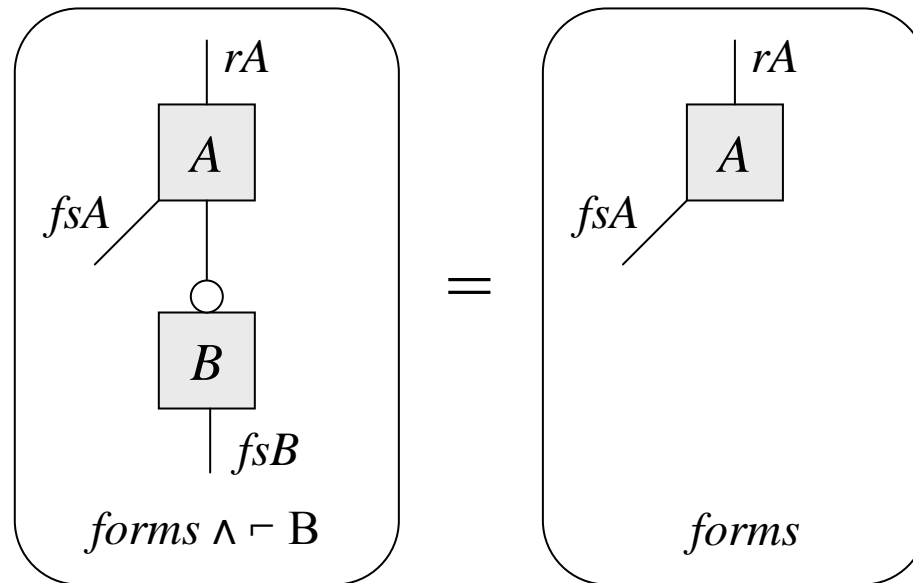
Replace Node (Law 4)



Remove Root (Law 5)



Remove Node (Law 6)



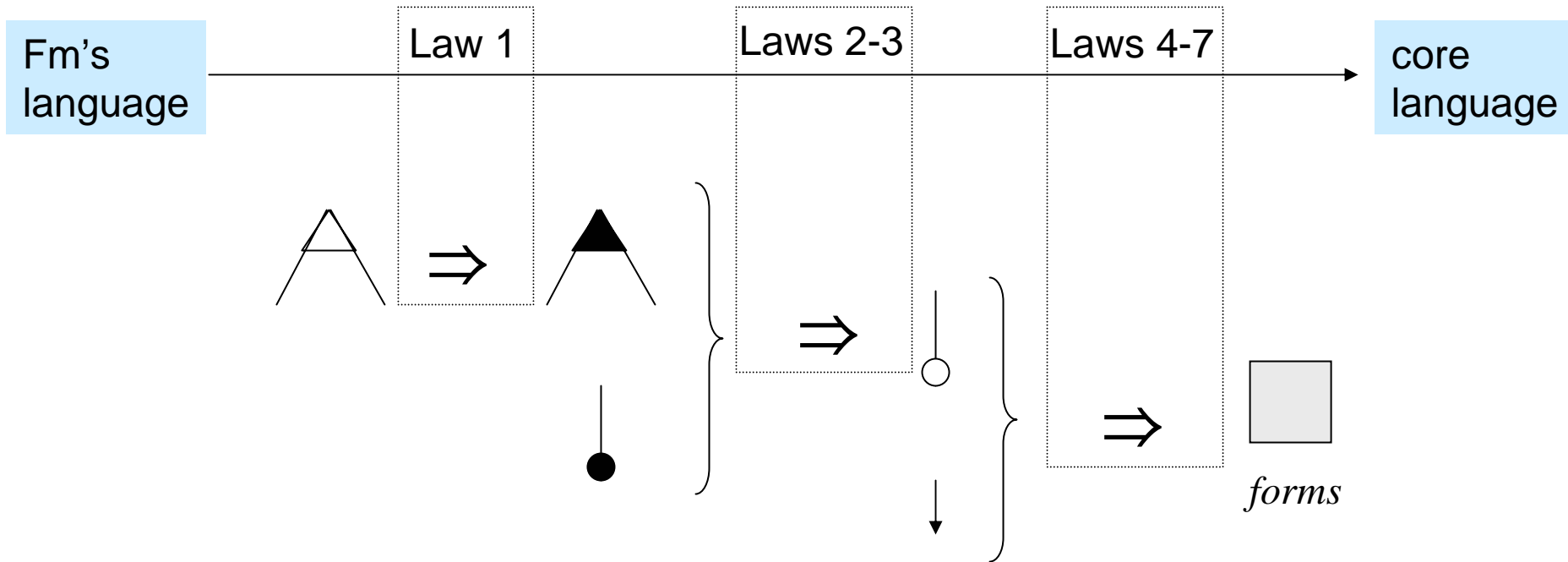
Add Formula (Law 7)

$$\boxed{\begin{array}{c} fs \\ forms \end{array}} = \boxed{\begin{array}{c} fs \\ forms \wedge f \end{array}}$$

(\leftrightarrow) f can be deduced from $forms$ and fs .

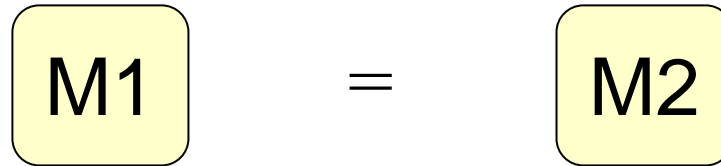
Proved in PVS

Reduction Strategy

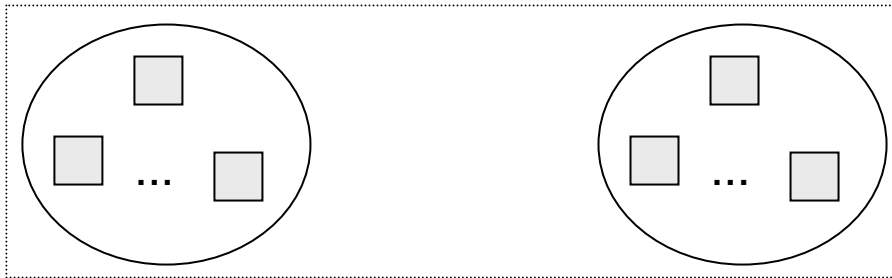


All laws are applied from left to right

Completeness (proof)



reduction
strategy
(Laws 1-6)



If they are equivalent, they
have the same features.

$$formsM1 = \dots = formsM2$$

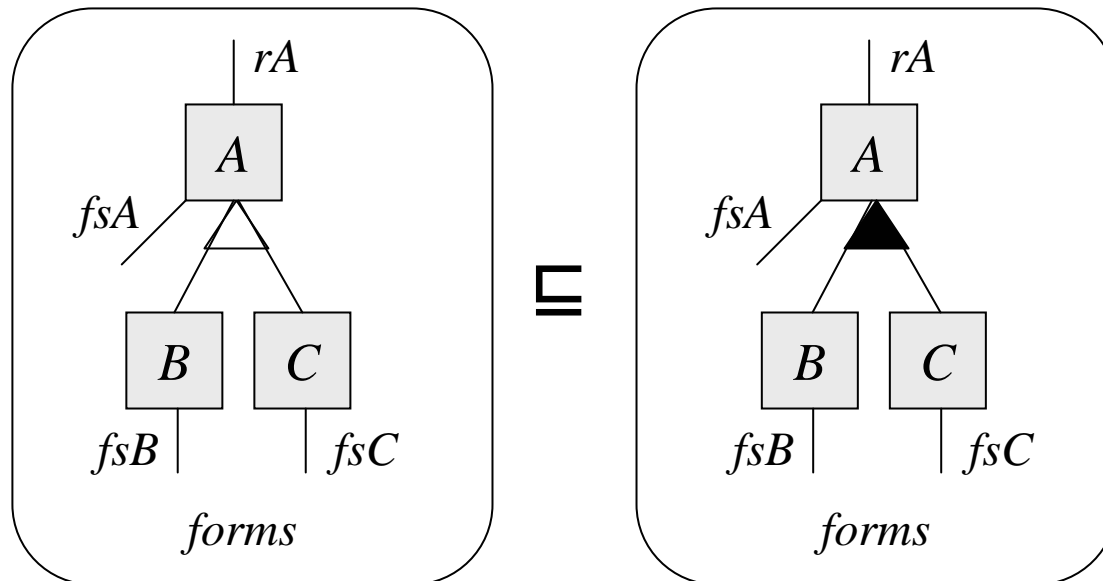
If they are equivalent, they have equivalent constraints. Since propositional logic's calculus is complete, we can always prove:
 $formsM1 = formsM2$. (Law 7)

Refinements

Refinements

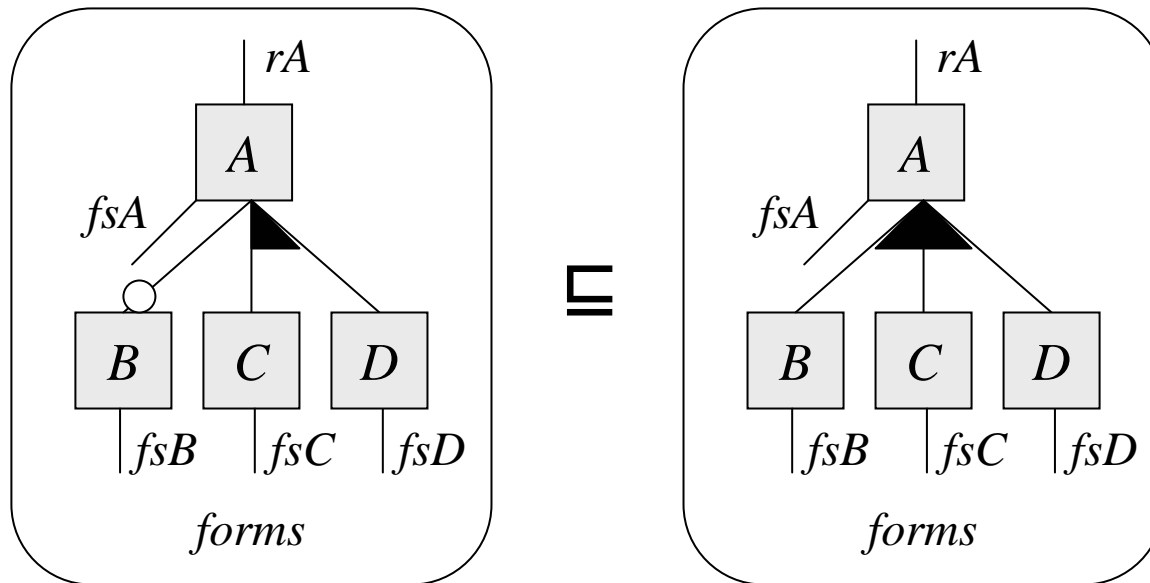
- We propose 12 refinements
- All of them are proved in PVS with respect to
 - static semantics
 - dynamic semantics
- It is important to mention that we have more transformations since most of them can be applied with more than two features

Change Alternative to Or (Ref 1)



Proved in PVS

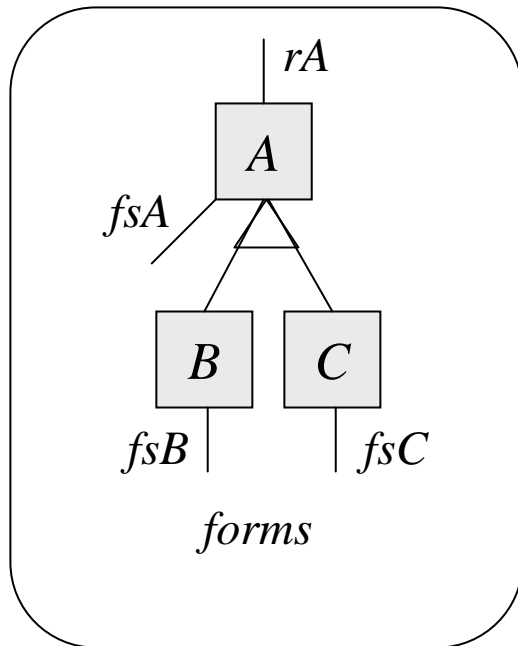
Collapse Optional and Or (Ref 2)



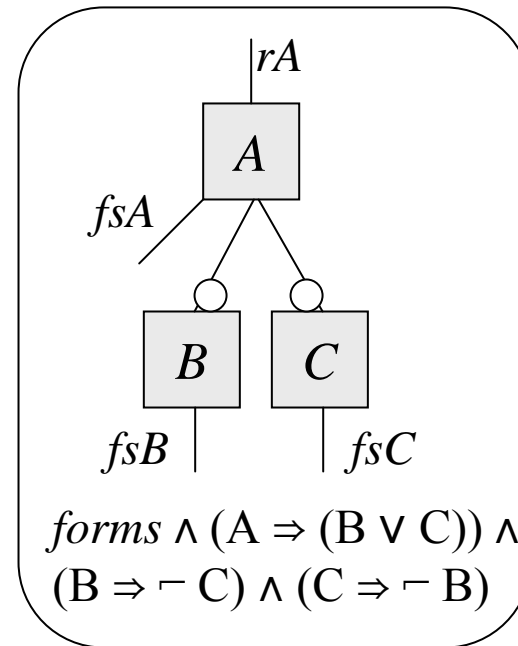
Proved in PVS

Change Alternative to Optional (Ref 3)

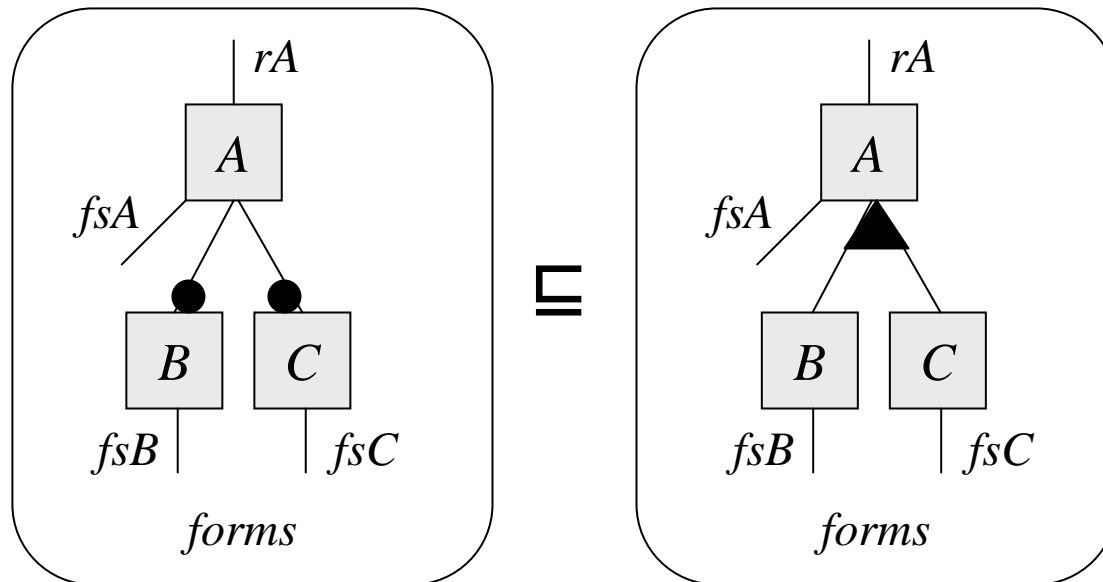
Law 1 (\rightarrow) \Rightarrow Law 2 (\rightarrow)



=

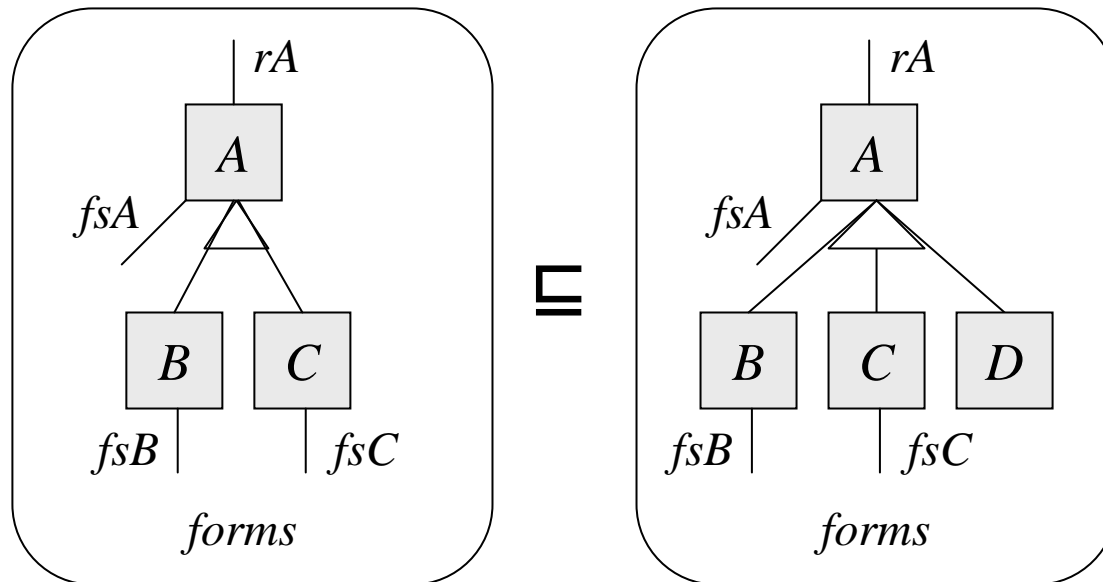


Add Or Between Mandatory (Ref 4)



Proved in PVS

Add New Alternative (Ref 5)



D is a new name

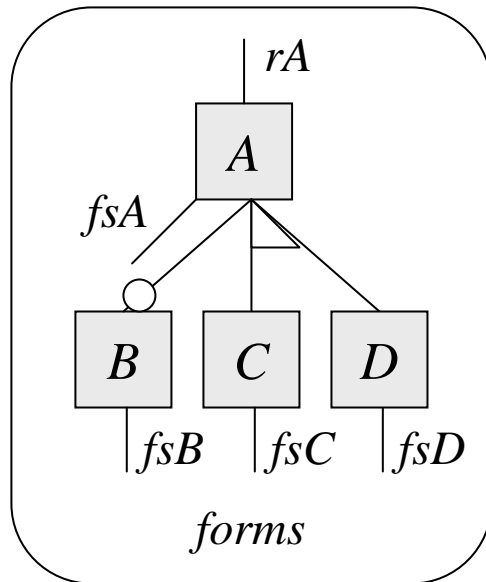
Proved in PVS

Collapse Optional and Alternative to Or (Ref 6)

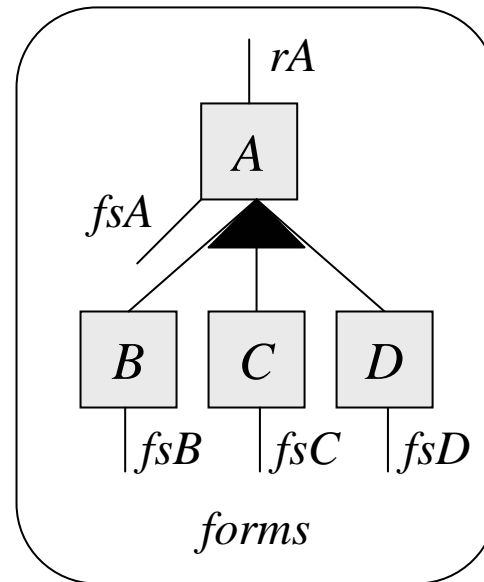
Ref 1

\Rightarrow

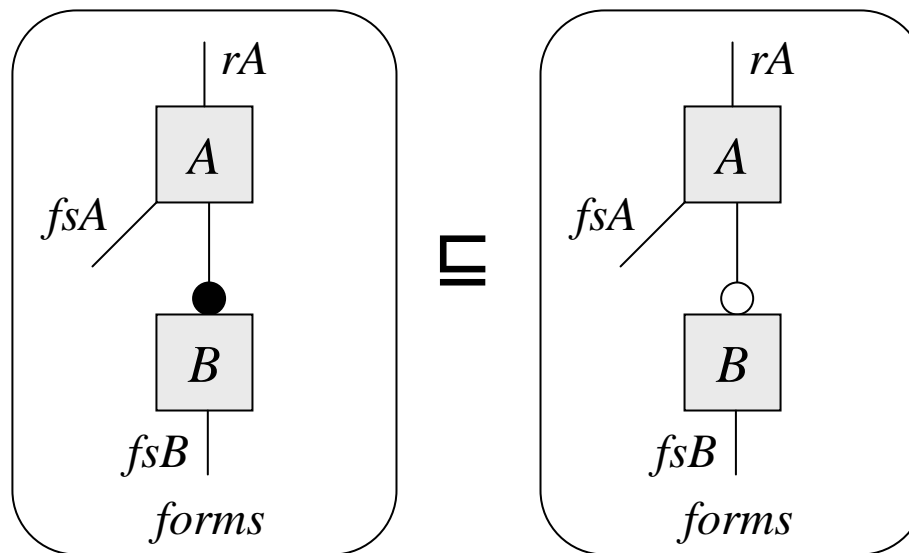
Ref 2



\sqsubseteq

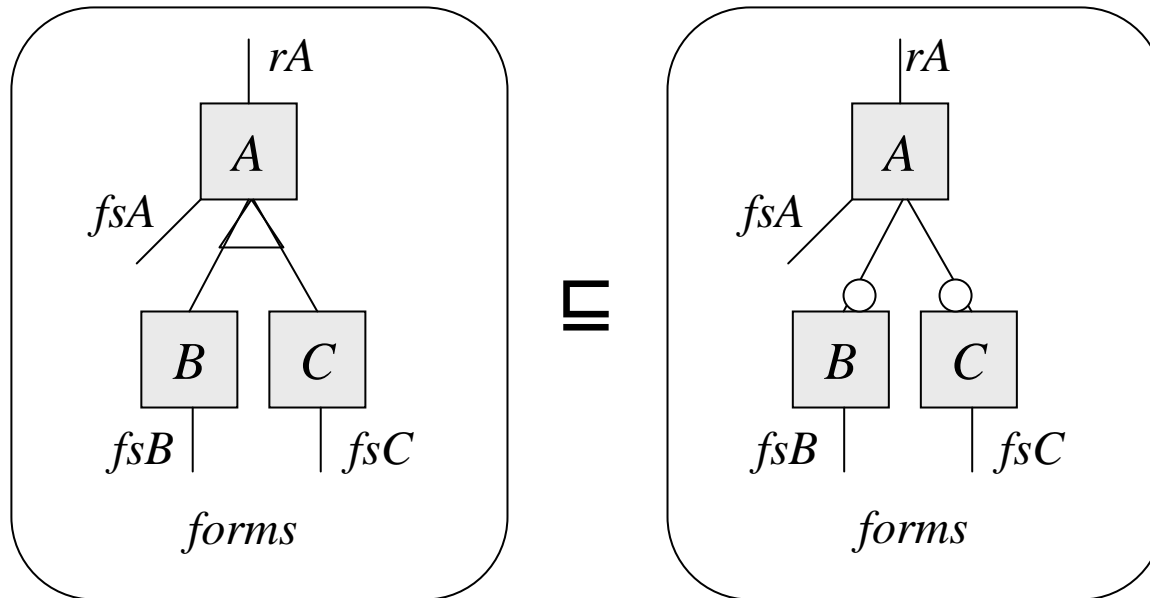


Remove Mandatory (Ref 7)



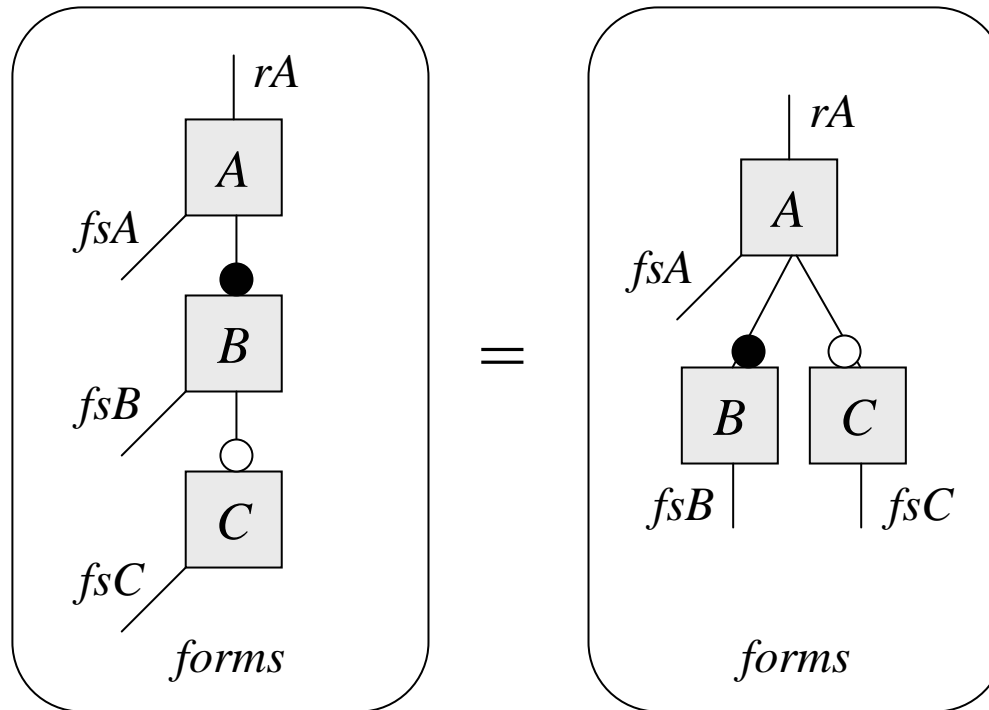
Proved in PVS

Remove Alternative (Ref 8)



Proved in PVS

Pull up Node (Ref 9-10)



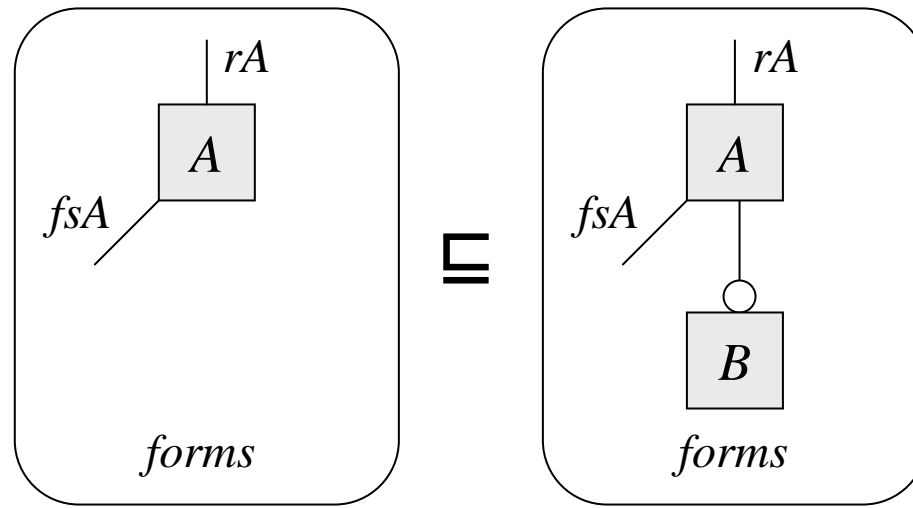
Proved in PVS

Remove Formula (Ref 11)

$$\boxed{\begin{array}{c} fs \\ forms \wedge f \end{array}} \sqsubseteq \boxed{\begin{array}{c} fs \\ forms \end{array}}$$

Proved in PVS

Add Optional Node (Ref 12)

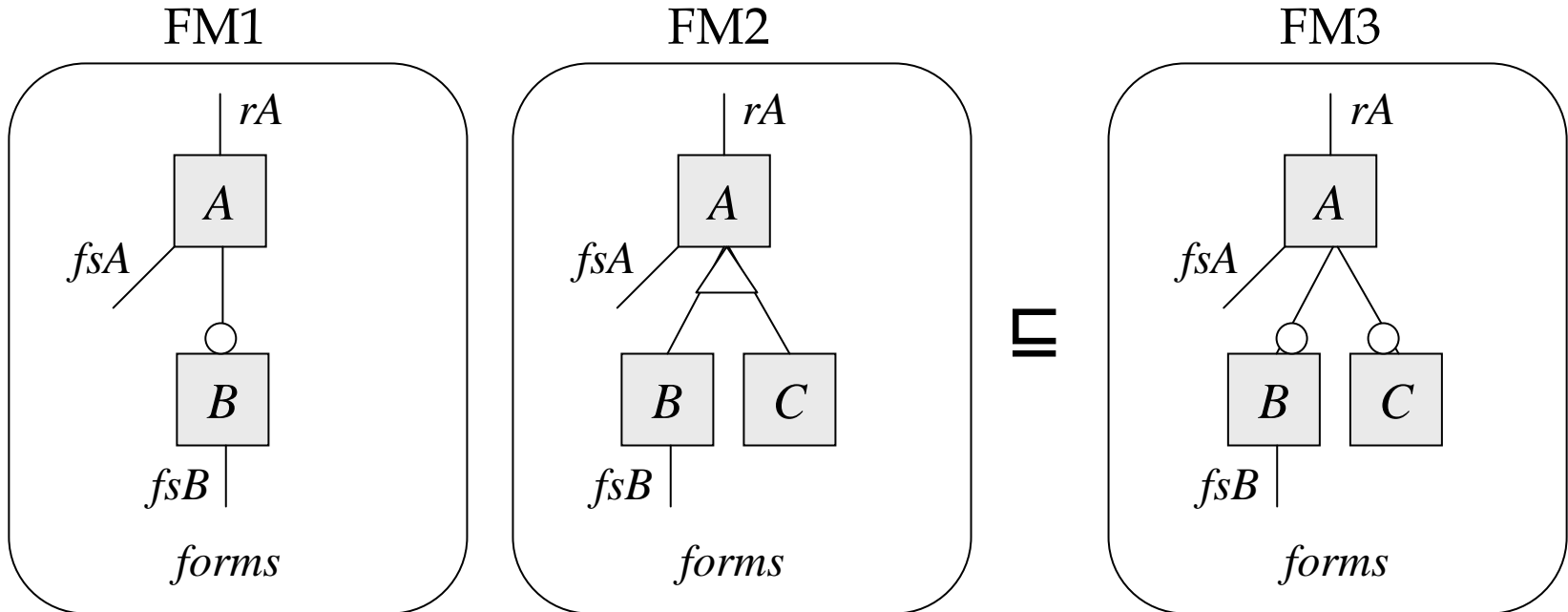


B is a new name

Proved in PVS

Extractive Refinements

Extractive 1

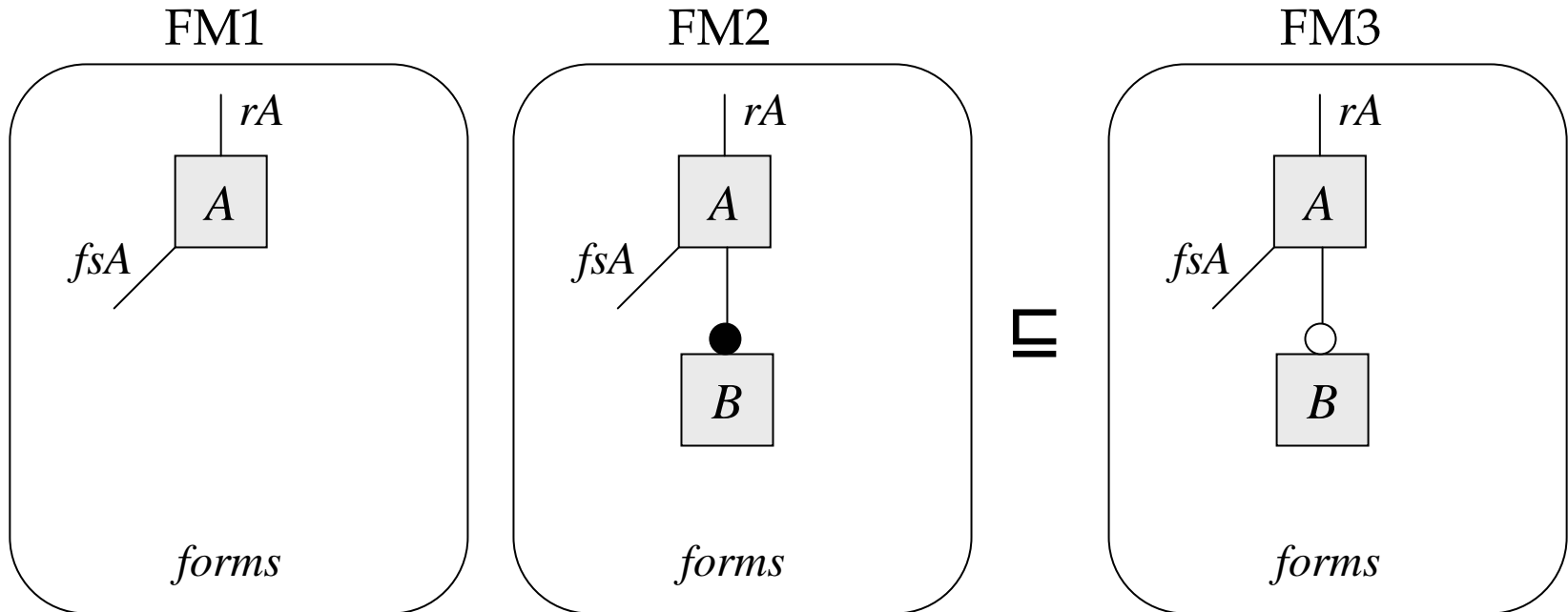


Proofs:

FM1 \sqsubseteq FM3: Ref 12 (intro C)

FM2 \sqsubseteq FM3: Ref 8

Extractive 2

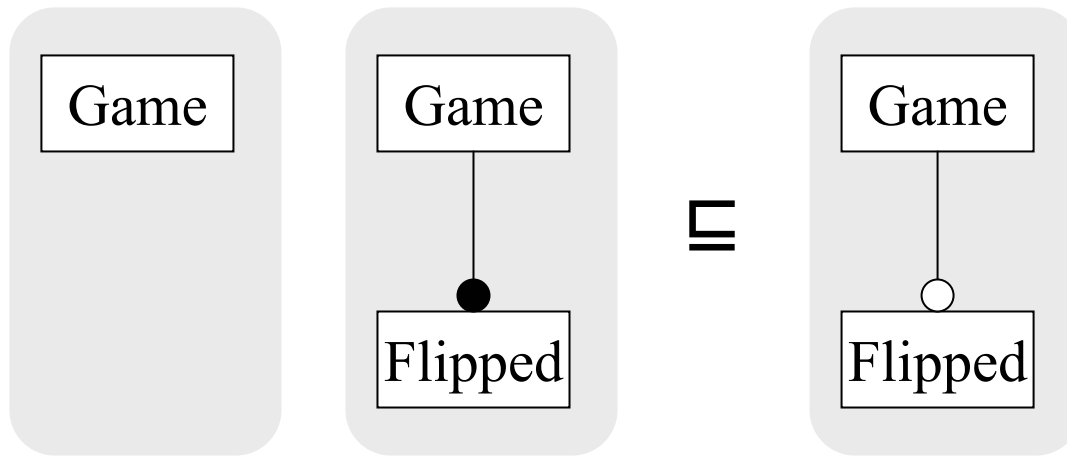


Proofs:

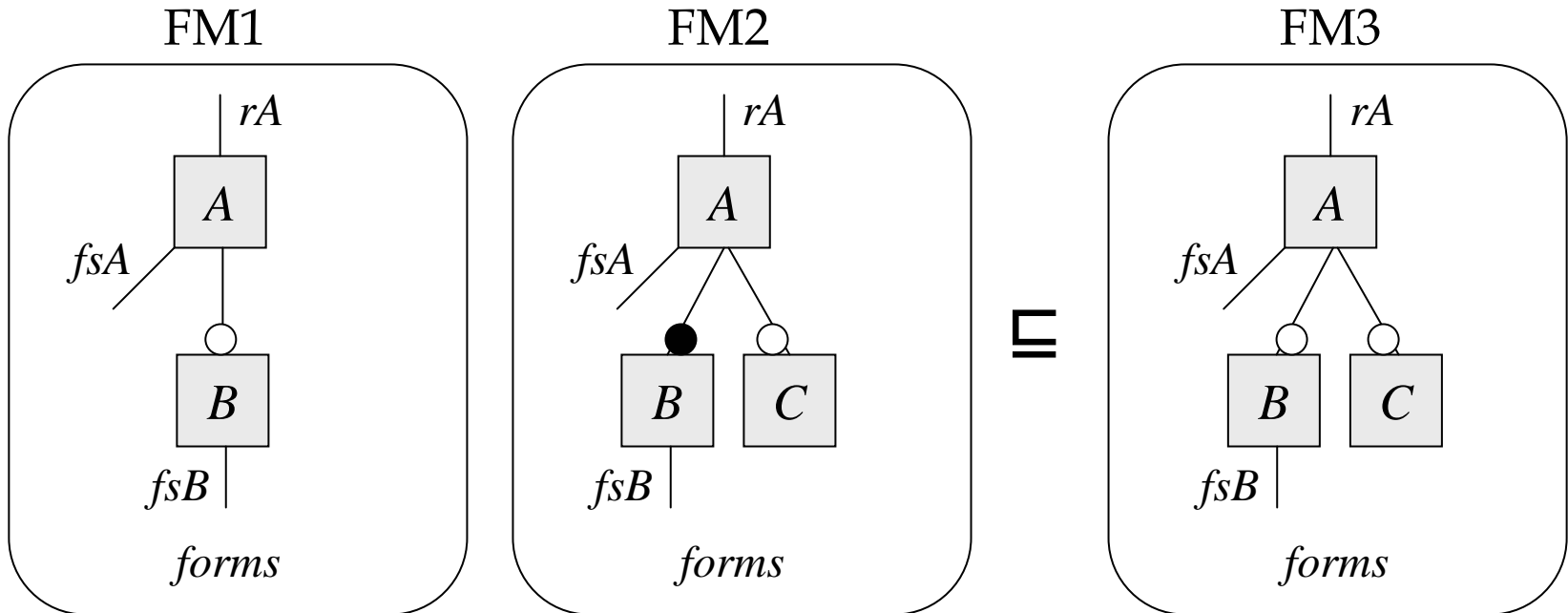
FM1 \sqsubseteq FM3: Ref 12 (intro B)

FM2 \sqsubseteq FM3: Ref 7 (B)

Example



Extractive 3

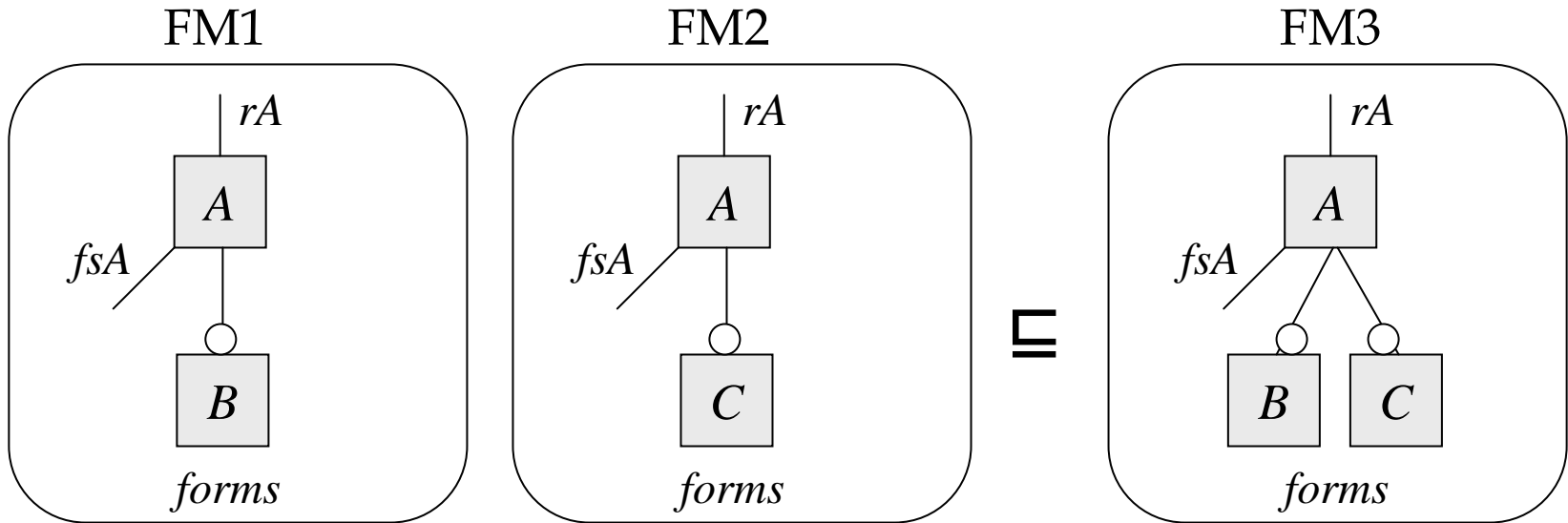


Proofs:

FM1 \sqsubseteq FM3: Ref 12 (intro C)

FM2 \sqsubseteq FM3: Ref 7 (B)

Extractive 4



Proofs:

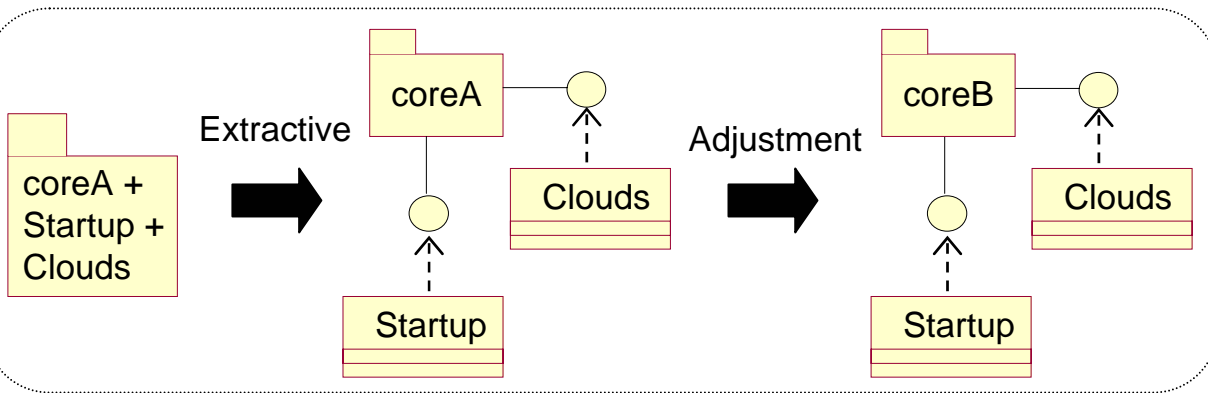
FM1 \sqsubseteq FM3: Ref 12 (intro C)

FM2 \sqsubseteq FM3: Ref 12 (intro B)

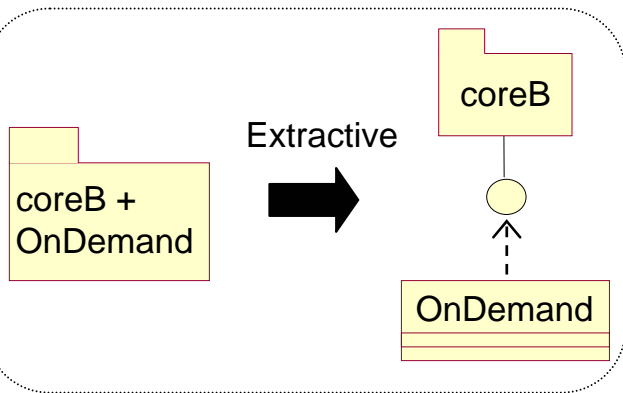
Evaluation

Case Study

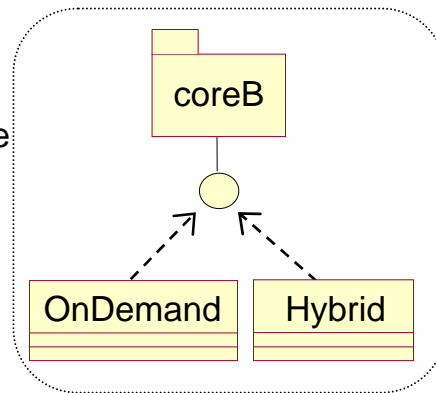
Product 1



Product 2



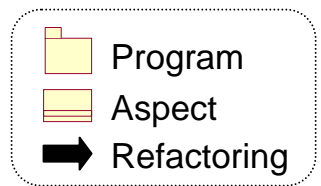
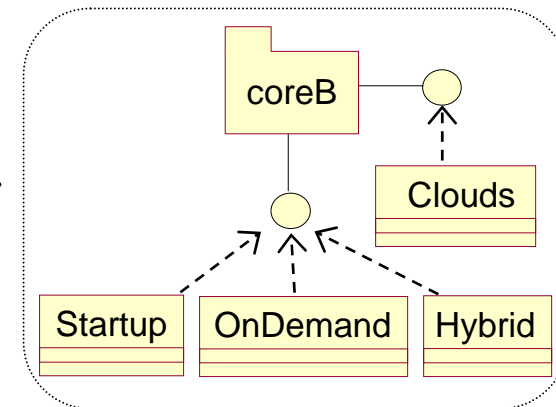
SPL 2



Extractive

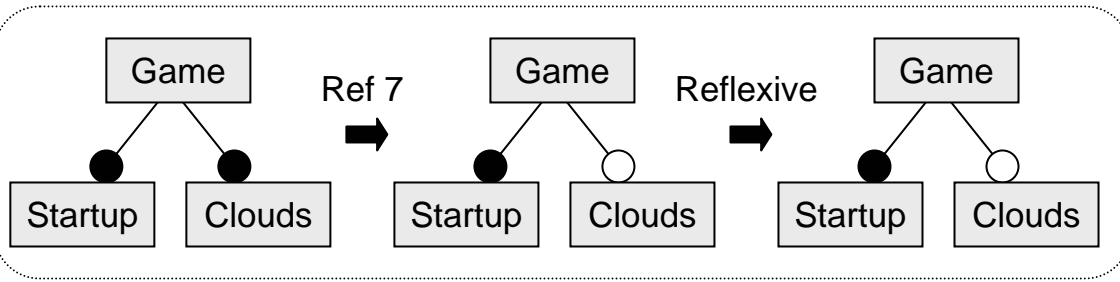


SPL 1-2



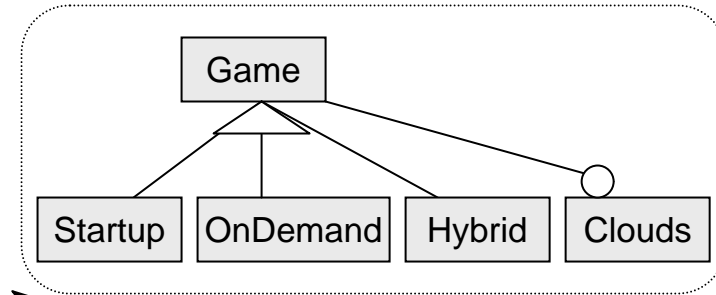
FM Reasoning - Case Study

Product 1

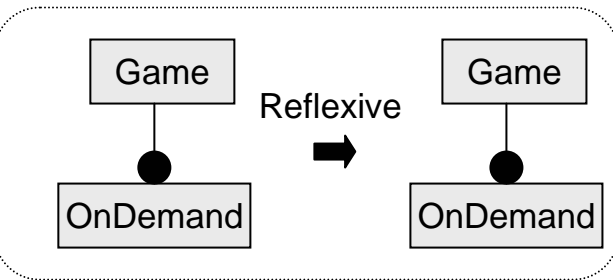


2x Ref 5

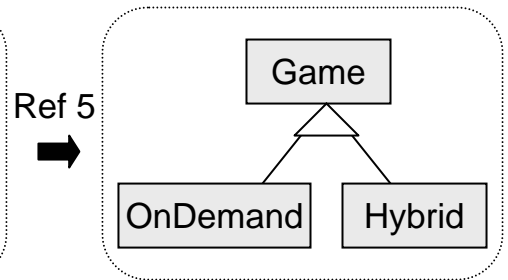
SPL 1-2



Product 2

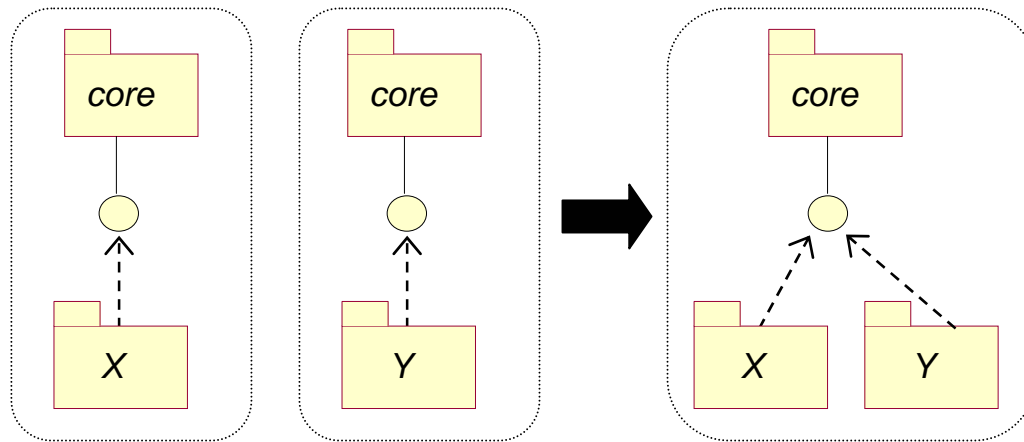


SPL 2



Ref 5, 12

Product Line Refactoring 1



X and Y do not have any classes, aspects, members and methods with the same name